

## Hamiltonian Chromatic Number of Graphs

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### Abstract

This paper studies the Hamiltonian coloring and Hamiltonian chromatic number for different graphs .the main results are1.For any integer n greater than or equal to three, Hamiltonian chromatic number of  $C_n$  is equal to n-2. 2. G is a graph obtained by adding a pendant edge to Hamiltonian graph H, and then Hamiltonian chromatic number of G is equal to n-1. 3. For every connected graph G of order n greater than or equal to 2, Hamiltonian chromatic number of G is not more than one increment of square of (n-2).

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**Key words:** Chromatic number, Hamiltonian coloring, Hamiltonian chromatic number, pendant edge, spanning connected graph.

### I. Introduction

Generally in a  $(d - 1)$  radio coloring of a connected graph G of diameter d, the colors assigned to adjacent vertices must differ by at least d-1, the colors assigned to two vertices whose distance is 2 must differ by at least d-2, and so on up to antipodal vertices, whose colors are permitted to be the same. For this reason,  $(d-1)$  radio colorings are also referred to as antipodal colorings.

In the case of an antipodal coloring of the path  $P_n$  of order  $n \geq 2$ , only the two end-vertices are permitted to be colored the same. If u and v are distinct vertices of  $P_n$  and  $d(u, v) = i$ , then  $|c(u)-c(v)| \geq n - 1 - i$ . Since  $P_n$  is a tree, not only is i the length of a shortest u - v path in  $P_n$ , it is the length of the only u - v path in  $P_n$ . In particular, is the length of a longest u - v path?

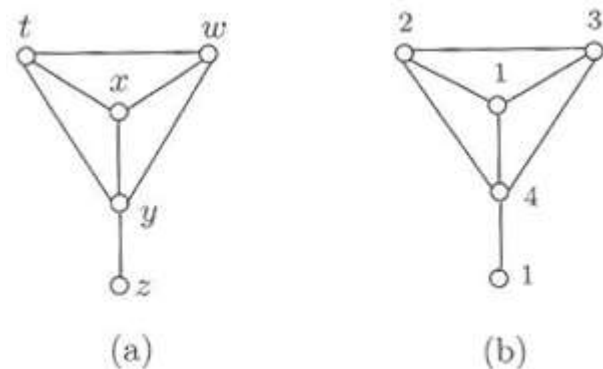
The detour distance  $D(u, v)$  between two vertices u and v in a connected graph G is defined as the length of a longest u - v path in G. Hence the length of a longest u - v path in  $P_n$  is  $D(u, v) = d(u, v)$ . Therefore, in the case of the path  $P_n$ , an antipodal coloring of  $P_n$  can also be defined as a vertex coloring c that satisfies.

$D(u,v) + |c(v)| \geq n - 1$ , for every two distinct vertices u and v of  $P_n$ .

**§1.1 Definition:** Vertex coloring c that satisfy were extended from paths of order n to arbitrary connected graphs of order n by Gary Chartrand, Ladislav Nebesky, and Ping Zhang .A **Hamiltonian coloring** of a connected graph G of order n is a vertex coloring c such that,  $D(u,v) + |c(v)| \geq n - 1$ , for every two

distinct vertices u and v of G. the largest color assigned to a vertex of G by c is called the **value** of c and is denoted by  $hc(c)$ . The **Hamiltonian chromatic number**  $hc(G)$  is the smallest value among all Hamiltonian colorings of G.

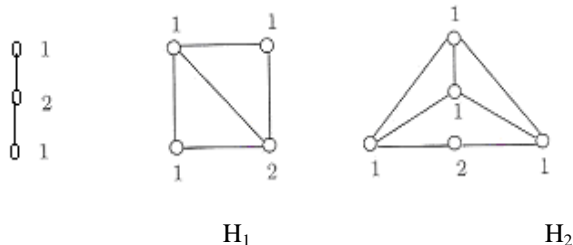
**EX:** Figure.1 (a) shows a graph H of order 5. A vertex coloring c of H is shown in Figure.1 (b). Since  $D(u, v) + |c(u) - c(v)| \geq 4$  for every two distinct vertices u and u of H, it follows that c is a Hamiltonian coloring and so  $hc(c) = 4$ . Hence  $hc(H) \geq 4$ . Because no two of the vertices t, w, x, and y are connected by a Hamiltonian path, these must be assigned distinct colors and so  $hc(H) \geq 4$ . Thus  $hc(H) = 4$ .



1. A graph with Hamiltonian chromatic number 4

If a connected graph G of order n has Hamiltonian chromatic number 1, then  $D(u,v) = n - 1$  for every two distinct vertices u and v of G and consequently G is Hamiltonian-connected, that is, every two vertices of G are connected by a

Hamiltonian path. Indeed,  $hc(G) = 1$  if and only if  $G$  is Hamiltonian – connected. Therefore, the Hamiltonian Chromatic Number of a connected graph  $G$  can be considered as a measure of how close  $G$  is to being Hamiltonian connected. That is the closer  $hc(G)$  is to 1, the closer  $G$  is to being Hamiltonian connected. The three graphs  $H_1$ ,  $H_2$  and  $H_3$  shown in below figure 2 are all close (in this sense) to being Hamiltonian-connected since  $hc(H_i) = 2$  for  $i = 1, 2, 3$ .



$H_3$   $H_1$   $H_2$   
 2. Three graphs with Hamiltonian chromatic number 2

**II. Theorem: For every integer  $n \geq 3$ ,  $hc(K_{1,n-1}) = (n-2)^2 + 1$**

Proof: Since  $hc(K_{1,2}) = 2$  (See  $H_1$  in Figure 2) we may assume that  $n \geq 4$ .

Let  $G = K_{1,n-1}$  where  $V(G) = \{v_1, v_2, \dots, v_n\}$  and  $v_n$  is the central vertex. Define the coloring  $c$  of  $G$  by  $c(v_n) = 1$  and  $C(v_i) = (n-1) + (i-1)(n-3)$  for  $1 \leq i \leq n-1$ . Then  $c$  is a Hamiltonian coloring of  $G$  and  $hc(G) \leq hc(c) = c(v_{n-1}) = (n-1) + (n-2)(n-3) = (n-2)^2 + 1$ . It remains to show that  $hc(G) \geq (n-2)^2 + 1$ .

Let  $c$  be a Hamiltonian coloring of  $G$  such that  $hc(c) = hc(G)$ . Because  $G$  contains no Hamiltonian path,  $c$  assigns distinct colors to the vertices of  $G$ . We may assume that  $C(v_1) < c(v_2) < \dots < c(v_{n-1})$ . We now consider three cases, depending on the color assigned to the central vertex  $v_n$ .

**Case 1.**

$c(v_n) = 1$ .

Since

$$D(v_1, v_n) = 1 \text{ and } D(v_i, v_{i+1}) = 2 \text{ for } 1 \leq i \leq n-2.$$

It follows that

$$C(v_{n-1}) \geq 1 + (n-2) + (n-2)(n-3) = (n-2)^2 + 1$$

And so

$$hc(G) = hc(c) = c(v_{n-1}) \geq (n-2)^2 + 1.$$

**Case 2.**

$$C(v_n) = hc(c)$$

Thus, in this case,

$$1 = c(v_1) < c(v_2) < \dots < c(v_{n-1}) < c(v_n)$$

Hence

$$C(v_n) \geq 1 + (n-2)(n-3) + (n-2) = (n-2)^2 + 1$$

And so

$$hc(G) = hc(c) = c(v_n) \geq (n-2)^2 + 1.$$

**Case 3.**

$C(v_j) < c(v_n) < c(v_{j+1})$  for some integer  $j$  with  $1 \leq j \leq n-2$ .

Thus  $c(v_1) = 1$  and  $c(v_{n-1}) = hc(c)$ .

in this case

$$C(v_j) \geq 1 + (j-1)(n-3),$$

$$C(v_n) \geq c(v_j) + (n-2)$$

$$C(v_{j+1}) \geq c(v_n) + (n-2) \text{ and}$$

$$C(v_{n-1}) \geq c(v_{j+1}) + [(n-1) - (j+1)](n-3).$$

Therefore,

$$C(v_{n-1}) \geq 1 + (j-1)(n-3) + 2(n-2) + (n-j-2)(n-3)$$

$$= (2n-3) + (n-3)^2 = (n-2)^2 + 2 > (n-2)^2 + 1.$$

$$\text{And so } hc(G) = hc(c) = c(v_{n-1}) > (n-2)^2 + 1.$$

Hence in any case,

$$hc(G) \geq (n-2)^2 + 1 \text{ and so } hc(G) = (n-2)^2 + 1.$$

**III. Theorem: For every integer  $n \geq 3$ ,  $hc(C_n) = n-2$ .**

**Proof.**

Since we noted that  $hc(C_n) = n-2$  for  $n = 3, 4, 5$ .

We may assume that  $n \geq 6$ . Let  $C_n = (v_1, v_2, \dots, v_n, v_1)$ . Because the vertex coloring  $c$  of  $C_n$  defined by  $c(v_1) = c(v_2) = 1$ ,  $c(v_{n-1}) = c(v_n) = n-2$  and  $c(v_i) = i-1$  for  $3 \leq i \leq n-2$  is a Hamiltonian coloring, it follows that  $hc(C_n) \leq n-2$ . Assume, to the contrary, that  $hc(C_n) < n-2$  for some integer  $n \geq 6$ . Then there exists a Hamiltonian  $(n-3)$  coloring  $c$  of  $C_n$ . We consider two cases, according to whether  $n$  is odd or  $n$  is even.

**Case 1.**

**n is odd:** Then  $n = 2k + 1$  for some integer  $k \geq 3$ . Hence there exists a Hamiltonian  $(2k-2)$  coloring  $c$  of  $C_n$ . Let,

$$A = \{1, 2, \dots, k-1\} \text{ and } B = \{k, k+1, \dots, 2k-2\}$$

For every vertex  $u$  of  $C_n$ , there are two vertices  $v$  of  $C_n$  such that  $D(u,v)$  is minimum (and  $d(u,v)$  is maximum), namely  $D(u,v) = d(u,v) + 1 = k+1$ . For  $u = v_i$ , these two vertices  $v$  are  $v_{i+k}$  and  $v_{i+k+1}$  (where the addition in  $i+k$  and  $i+k+1$  is performed modulo  $n$ ). Since  $c$  is a Hamiltonian coloring.

$$D(u, v) + |c(u) - c(v)| \geq n-1 = 2k. \text{ Because } D(u, v) = k+1, \text{ it follows that}$$

$$|c(u) - c(v)| \geq k-1.$$

Therefore, if  $c(u) \in A$ , then the colors of these two vertices  $v$  with this property must belong to  $B$ . In particular, if  $c(v_i) \in A$ , then  $(v_{i+k}) \in B$ . Suppose that there are  $a$  vertices of  $C_n$  whose colors belong to  $A$  and  $b$  vertices of  $C_n$  whose colors belong to  $B$ . Then  $b \geq a$ . However, if  $c(v_i) \in B$ , then  $c(v_{i+k})$  belongs to  $A$  implying that  $a \geq b$  and so,  $a=b$ . since  $a+b=n$  and  $n$  is odd, this is impossible.

**Case 2.**

**n is even:** Then  $n = 2k$  for some integer  $k \geq 3$ . Hence there exists a Hamiltonian  $(2k-3)$  coloring  $c$  of  $C_n$ . For every vertex  $u$  of  $C_n$ , there is a unique vertex  $v$  of  $C_n$  for which  $D(u, v)$  is minimum (and  $d(u, v)$  is maximum), namely,  $d(u, v) = k$ . For  $u = v_i$ , this

vertex  $v$  is  $v_{i+k}$  (where the addition in  $i + k$  is performed modulo  $n$ ).

Since  $c$  is a Hamiltonian coloring,  $D(u,v) + |c(u) - c(v)| \geq n - 1 = 2k - 1$ . Because  $D(u, v) = k$ , it follows that  $|c(u) - c(v)| \geq k - 1$ . This implies, however, that if

$c(u) = k-1$ , then there is no color that can be assigned to  $u$  to satisfy this requirement. Hence no vertex of  $C_n$  can be assigned the color  $k-1$  by  $c$ .

Let,  $A = \{1, 2, \dots, k-2\}$  and  $B = \{k, k+1, \dots, 2k-3\}$ .

Thus  $|A| = |B| = k - 2$ . If  $c(v_i) \in A$ , then  $c(v_{i+k}) \in B$ . Also, if  $c(v_i) \in B$ , then  $c(v_{i+k}) \in A$ . Hence there are  $k$  vertices of  $C_n$  assigned colors from  $B$ . Consider two adjacent vertices of  $C_n$ , one of which is assigned a color from  $A$  and the other is assigned a color from  $B$ . We may assume that  $c(v_1) \in A$  and  $c(v_2) \in B$ . Then  $c(v_{k+1}) \in B$ . Since  $D(v_2, v_{k+1}) = k+1$ , it follows that  $|c(v_2) - c(v_{k+1})| \geq k - 2$ . Because  $c(v_2) \in B$ ,  $c(v_{k+1}) \in B$ , this implies that one of  $c(v_2)$  and  $c(v_{k+1})$  is at least  $2k-2$ . This is a contradiction.

**§ 3.1 Proposition: If  $H$  is a spanning connected sub graph of a graph  $G$ , then  $hc(G) \leq hc(H)$**

**Proof.**

Suppose that the order of  $H$  is  $n$ . Let  $c$  be a Hamiltonian coloring of  $H$  such that  $hc(c) = hc(H)$ . Then  $D_H(u,v) + |c(u) - c(v)| \geq n - 1$  for every two distinct vertices  $u$  and  $v$  of  $H$ . since  $D_G(u,v) \geq D_H(u,v)$  for every two distinct vertices  $u$  and  $v$  of  $H$ , it follows that  $D_G(u,v) + |c(u) - c(v)| \geq n - 1$  and so  $c$  is a Hamiltonian coloring of  $G$  as well. Hence  $hc(G) \leq hc(c) = hc(H)$ .

**§ 3.2 Proposition: Let  $H$  be a Hamiltonian graph of order  $n - 1 \geq 3$ . If  $G$  is a graph obtained by adding a pendant edge to  $H$ , then  $hc(G) = n - 1$ .**

**Proof.** Suppose that  $C = (v_1, v_2, \dots, v_{n-1}, v_1)$  is a Hamiltonian cycle of  $H$  and  $v_1 v_n$  is the pendant edge of  $G$ . Let  $c$  be a Hamiltonian coloring of  $G$ . Since  $D_G(u, v) \leq n-2$  for every two distinct vertices  $u$  and  $v$  of  $C$ , no two vertices of  $C$  can be assigned the same color by  $c$ . Consequently,  $hc(c) > n - 1$  and so  $hc(G) \geq n - 1$ .

Now define a coloring  $c'$  of  $G$  by

$$c'(v_i) = \begin{cases} i & \text{if } 1 < i < n - 1 \\ n - 1 & \text{if } i = n. \end{cases}$$

We claim that  $c'$  is a Hamiltonian coloring of  $G$ . First let  $v_j$  and  $v_k$  be two vertices of  $C$  where  $1 \leq j < k \leq n - 1$ . The  $|c'(v_j) - c'(v_k)| = k - j$  and

$$D(v_j, v_k) = \max \{k-j, (n-1) - (k-j)\}.$$

In either case,  $D(v_j, v_k) \geq n-1 + j - k$  and so

$$D(v_j, v_k) + |c'(v_j) - c'(v_k)| \geq n-1.$$

For  $1 \leq j \leq n-1$ ,  $|c'(v_j) - c'(v_n)| = n-1-j$ , while

$$D(v_j, v_n) \geq \max \{j, n-j+1\}$$

And so,  $D(v_j, v_n) \geq j$ .

Therefore,

$$D(v_j, v_n) + |c'(v_j) - c'(v_n)| \geq n-1.$$

Hence, as claimed,  $c'$  is a Hamiltonian coloring of  $G$  and so  $hc(G) \leq hc(c') = c'(v_n) = n-1$ .

**IV. Theorem: for every connected graph  $G$  of order  $n \geq 2$ ,  $hc(G) \leq (n-2)^2 + 1$ .**

**Proof.** First, if  $G$  contains a vertex of degree  $n-1$ , then  $G$  contains the star  $K_{1, n-1}$  as a spanning sub graph. Since  $hc(K_{1, n-1}) = (n-2)^2 + 1$  it follows by proposition 1 that  $hc(G) \leq (n-2)^2 + 1$ . Hence we may assume that  $G$  contains a spanning tree  $T$  that is not a star and so its complement  $T$  contains a Hamiltonian path  $P = (v_1, v_2, \dots, v_n)$ . Thus  $v_i v_{i+1} \notin E(T)$  for  $1 \leq i \leq n-1$  and so  $D_T(v_i, v_{i+1}) \geq 2$ . Define a vertex coloring  $c$  of  $T$  by

$$C(v_i) = (n-2) + (i-2)(n-3) \text{ for } 1 \leq i \leq n.$$

Hence

$$hc(c) = c(v_n) = (n-2) + (n-2)(n-3) = (n-2)^2$$

Therefore, for integers  $i$  and  $j$  with  $1 \leq i < j \leq n$ ,

$$|c(v_i) - c(v_j)| = (j-i)(n-3).$$

If  $j = i+1$ , then

$$D(v_i, v_j) + |c(v_i) - c(v_j)| \geq 1 + 2(n-3) = 2n-5 \geq n-1.$$

Thus  $c$  is a Hamiltonian coloring of  $T$ . therefore,

$$hc(G) \leq hc(T) \leq hc(c) = c(v_n) = (n-2)^2 < (n-2)^2 + 1,$$

Which completes the proof

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